

Inst.: Data- och Informationsteknik  
Kursnamn: Logic in Computer Science  
Examinator: Thierry Coquand  
Kurs: DIT200/DAT060

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No help documents

Telefonvakt: akn. 1030

*All answers and solutions must be carefully motivated!*

total 30;  $\geq 14$ : 3,  $\geq 19$ : 4,  $\geq 25$ : 5

total 30;  $\geq 14$ : G,  $\geq 21$ : VG

1. Give a proof in natural deduction of
  - a)  $((p \rightarrow q) \wedge (p \rightarrow r)) \rightarrow (p \rightarrow (q \wedge r))$  (2p)
  - b)  $(\exists x.(P(x) \vee Q(x))) \rightarrow ((\exists x.P(x)) \vee (\exists x.Q(x)))$  (2p)
  - c)  $(\forall x.\forall y.(S(y) \rightarrow F(x))) \rightarrow ((\exists y.S(y)) \rightarrow \forall x.F(x))$  (2p)
  
2. Say whether these consequences are true or false (if it is true give a proof, and if it is false give a model where it is not valid):
  - a)  $(\exists x.(P(x) \rightarrow Q(x))) \rightarrow ((\exists x.P(x)) \rightarrow (\exists y.Q(y)))$  (2p)
  - b)  $[(\forall x.\exists y.R(x, y)) \wedge (\forall y.\exists x.R(x, y))] \rightarrow \forall x.R(x, x)$  (2p)
  
3. We consider the following language: we have one unary predicate symbol  $P$  and one unary function symbol  $f$  and one constant  $a$ . Define what is a *model* for this language (2p). Give an example of a formula in this language, containing an occurrence of  $a$ ,  $f$  and  $P$ , which holds in some model but not in all models (1p).
  
4. Compute a CNF for the formula
 
$$(p \rightarrow q) \vee (q \wedge r)$$
 (1p)
 
$$(p \rightarrow (q \rightarrow r)) \rightarrow ((p \rightarrow q) \rightarrow r)$$
 (1p)
  
5. We consider the following CTL model  $(S, \rightarrow, L)$  where  $S$  has 4 states  $s_0, s_1, s_2, s_3$  and we have  $s_0 \rightarrow s_0, s_0 \rightarrow s_1, s_0 \rightarrow s_3, s_1 \rightarrow s_1, s_1 \rightarrow s_2, s_2 \rightarrow s_0, s_2 \rightarrow s_3, s_3 \rightarrow s_0$  and  $L(s_0) = \{a\}, L(s_1) = \{a, c\}, L(s_2) = \{b\}$  and  $L(s_3) = \{b, c\}$ . What are the states  $s$  for which we have?
  - a)  $s \models EX (EX c)$  (1p)?
  - b)  $s \models AG (EF b)$  (1p)
  
6. We consider the language with one unary predicate symbol  $P$  and one function symbol  $f$  and one constant  $a$ . We consider the formulae  $\phi_1 = P(a)$  and  $\phi_2 = \forall x.(P(x) \rightarrow P(f(x)))$ . Explain why we don't have  $\phi_1, \phi_2 \vdash \forall x.P(x)$  (2p). Give a model where we have  $\phi_1 \wedge \phi_2 \rightarrow \forall x.P(x)$  (2p).

Please turn the page!

7. Is the following LTL formula

$$F(\varphi) \rightarrow \varphi \vee X(X(F(\varphi)))$$

valid? Motivate your answer (2p)

8. Is the following entailment valid? (2p)

$$(p \rightarrow q) \rightarrow r, \neg r \wedge \neg s, (q \rightarrow p) \vee t, t \rightarrow (r \vee p) \vdash t$$

9. We consider the function  $F : X \mapsto (\{1, 3\} \cup X) \cap \{1, 4\}$  which takes as argument a subset to  $\{1, 2, 3, 4\}$  and return a subset to  $\{1, 2, 3, 4\}$ . What is the least fixed point of  $F$  and the greatest fixed point of  $F$  (2p)? Can you give another fixed point of  $F$  (1p)?

10. Explain why  $\rightarrow, \neg$  is a complete set of connectives (2p).

Good Luck!

Thierry and Jan